

Wet Chemical Etching in Semiconductor Fabrication

Physical problem A gap of width $2a$ and length L is to be etched in a flat plate. The remainder of the plate is covered with a protective (photoresist) layer. Since it is assumed that L is much larger than $2a$, the problem can be considered as two dimensional.

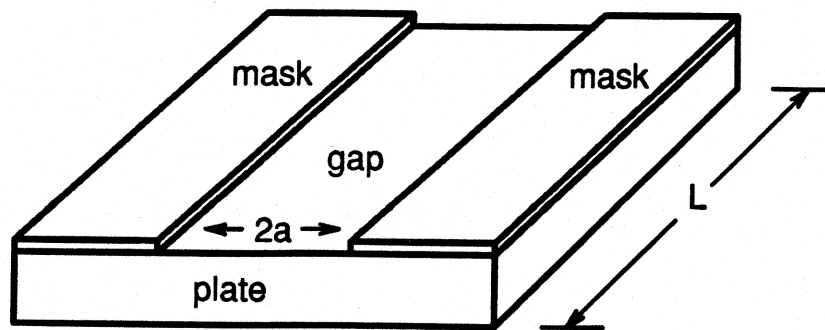


Figure 1: Physical Problem.

Assumptions:

- there is no convection in the etching medium;
- the etching process is isotropic;
- the thickness of the photoresist layer is infinitely small;
- only one component of the etching liquid determines the process.

Mathematical Model

The etching fluid $\Omega(t)$ is bounded by the outer boundary Γ_1 ; the photoresist layer $\Gamma_2(t)$, and the moving boundary $S(t)$. $D \setminus \Omega(t)$ denotes part of the solid.

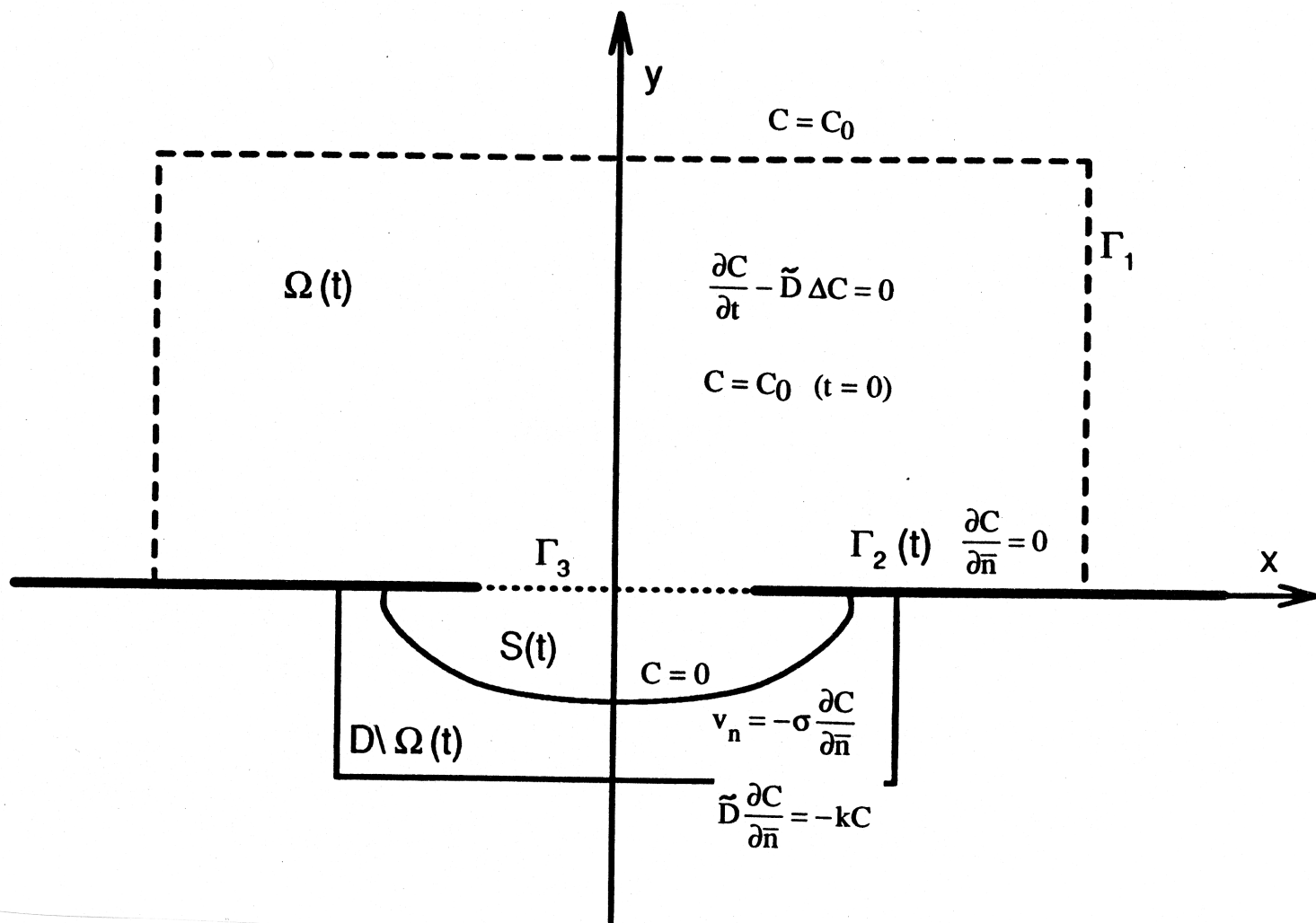


Figure 2: Side view of physical problem showing mathematical solution setup.

Fixed Domain Formulation

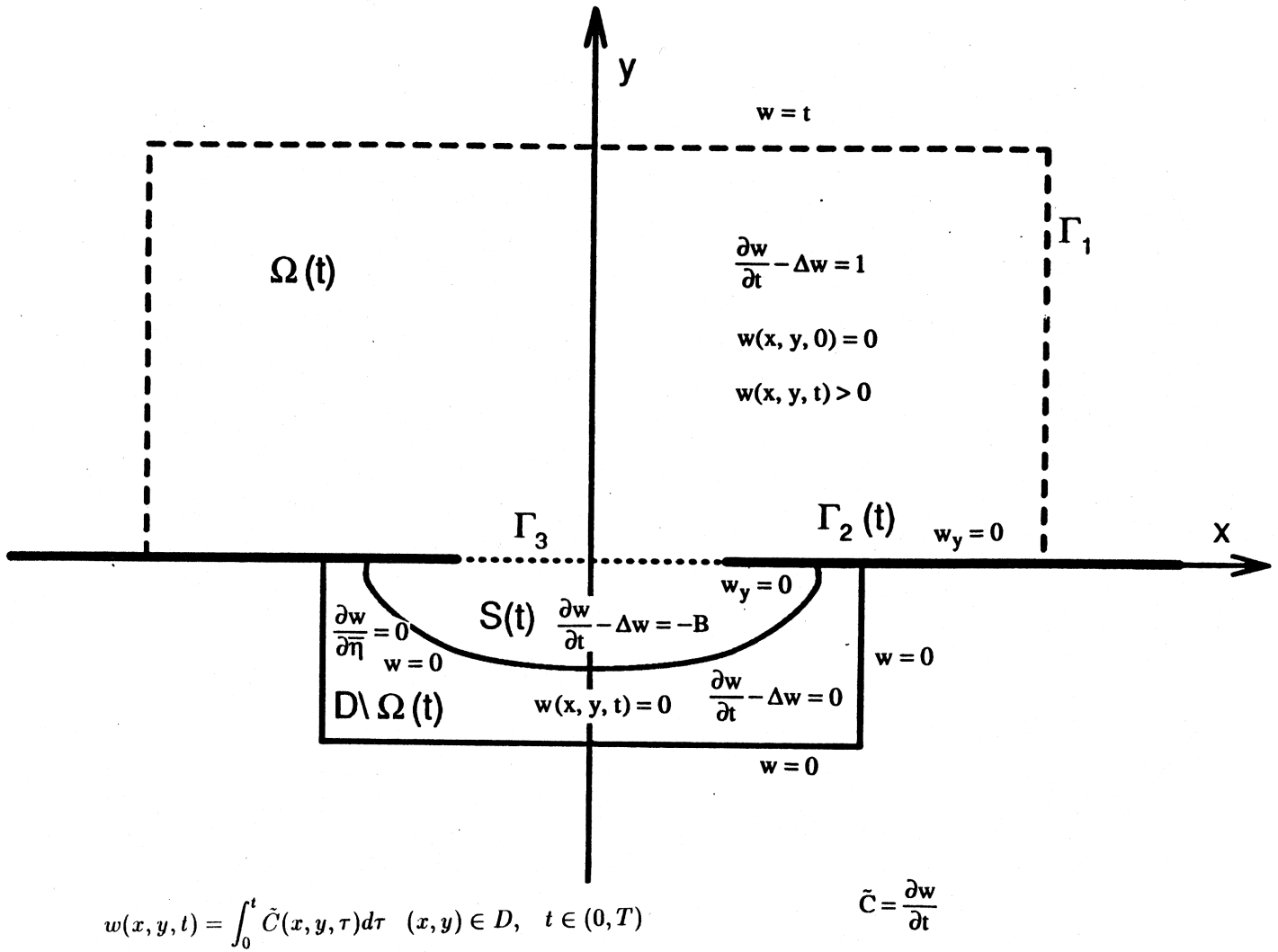


Figure 3: Fixed domain mathematical formulation.

Numerical Algorithm

The basic numerical algorithm is:

$$\begin{aligned}(w_1)_{i,j,k}^{(n+1/2)} = & \left(\frac{1}{\Delta t} + \frac{2}{(\Delta x)^2} + \frac{2}{(\Delta y)^2} \right)^{-1} \left\{ C_0 + \frac{1}{\Delta t} (w_1)_{i,j,k-1} \right. \\ & + \frac{1}{(\Delta x)^2} \left[(w_1)_{i-1,j,k}^{(n+1)} + (w_1)_{i+1,j,k}^{(n)} \right] \\ & \left. + \frac{1}{(\Delta y)^2} \left[(w_1)_{i,j-1,k}^{(n+1)} + (w_1)_{i,j+1,k}^{(n)} \right] \right\}\end{aligned}$$

with

$$(w_1)_{i,j,k}^{(n+1)} = (w_1)_{i,j,k}^{(n)} + \theta [(w_1)_{i,j,k}^{(n+1/2)} - (w_1)_{i,j,k}^{(n)}]$$

in $\Omega(0)$ and

$$\begin{aligned}(w_2)_{i,j,k}^{(n+1/2)} = & \left(\frac{1}{\Delta t} + \frac{32}{(\Delta x)^2} + \frac{32}{(\Delta y)^2} \right)^{-1} \left\{ -B + \frac{1}{\Delta t} (w_2)_{i,j,k-1} \right. \\ & + \frac{16}{(\Delta x)^2} \left[(w_2)_{i-1,j,k}^{(n+1)} + (w_2)_{i+1,j,k}^{(n)} \right] \\ & \left. + \frac{16}{(\Delta y)^2} \left[(w_2)_{i,j-1,k}^{(n+1)} + (w_2)_{i,j+1,k}^{(n)} \right] \right\}\end{aligned}$$

with

$$(w_2)_{i,j,k}^{(n+1)} = \max\{0, (w_2)_{i,j,k}^{(n)} + \theta [(w_2)_{i,j,k}^{(n+1/2)} - (w_2)_{i,j,k}^{(n)}]\}$$

in D_1

Domain Decomposition

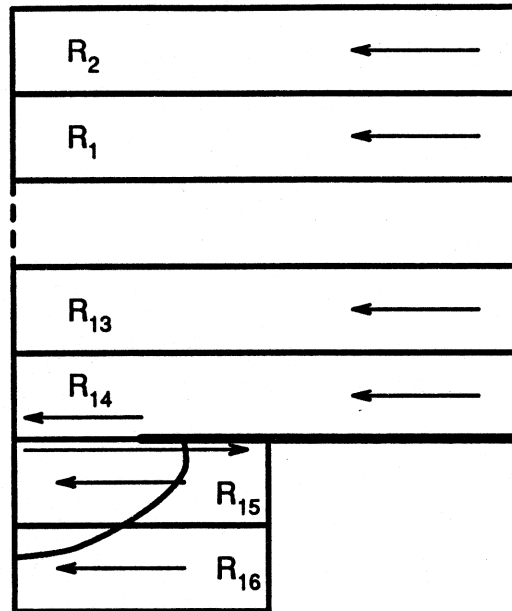


Figure 4: Domain decomposition of mathematical problem into sixteen subregions showing the flow of computations in each.

Load Balancing

Table: Load Balancing Information for Example

Nodes	2	4	8	16	32	64
bottom nodes	1	1	1	2	4	8
bottom points	12880	12880	12880	6440	3220	1610
top nodes	1	3	7	14	28	56
top points	89880	≤ 30174	12840	6420	3210	1605
diff points	77000	17294	40	20	10	5

Speedup

$$S_p = \frac{T_1}{T_p}$$

where T_p is the time required to execute the algorithm using p processors and T_1 is the time required to execute the same program on a single processor.

Efficiency

$$e_p = \frac{S_p}{p}$$

Results:

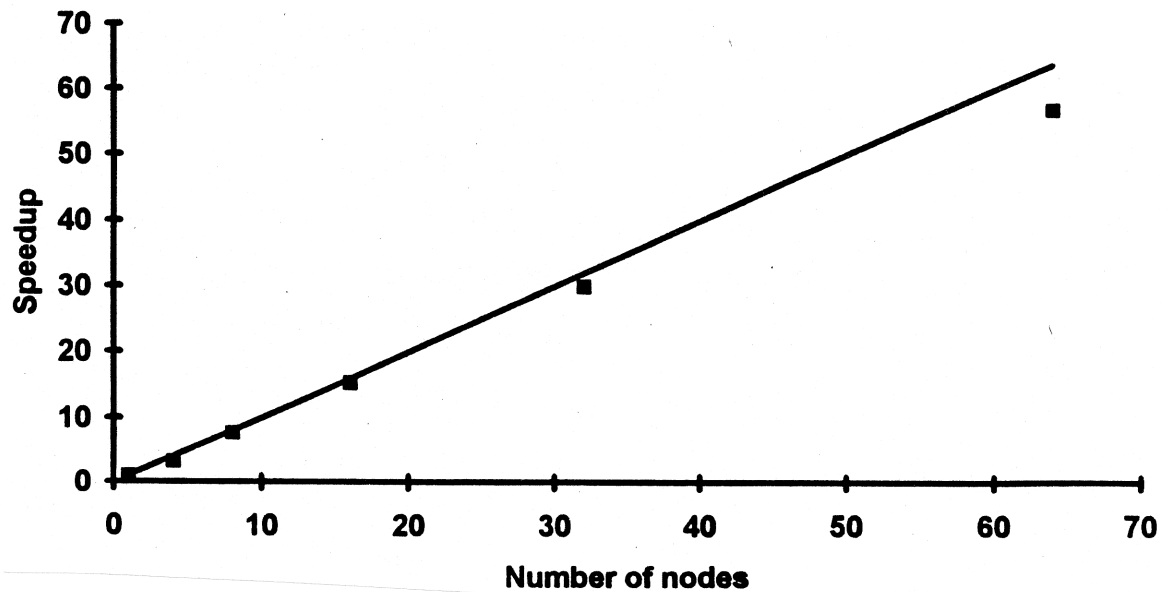


Figure 5: Speedups for various numbers of nodes.

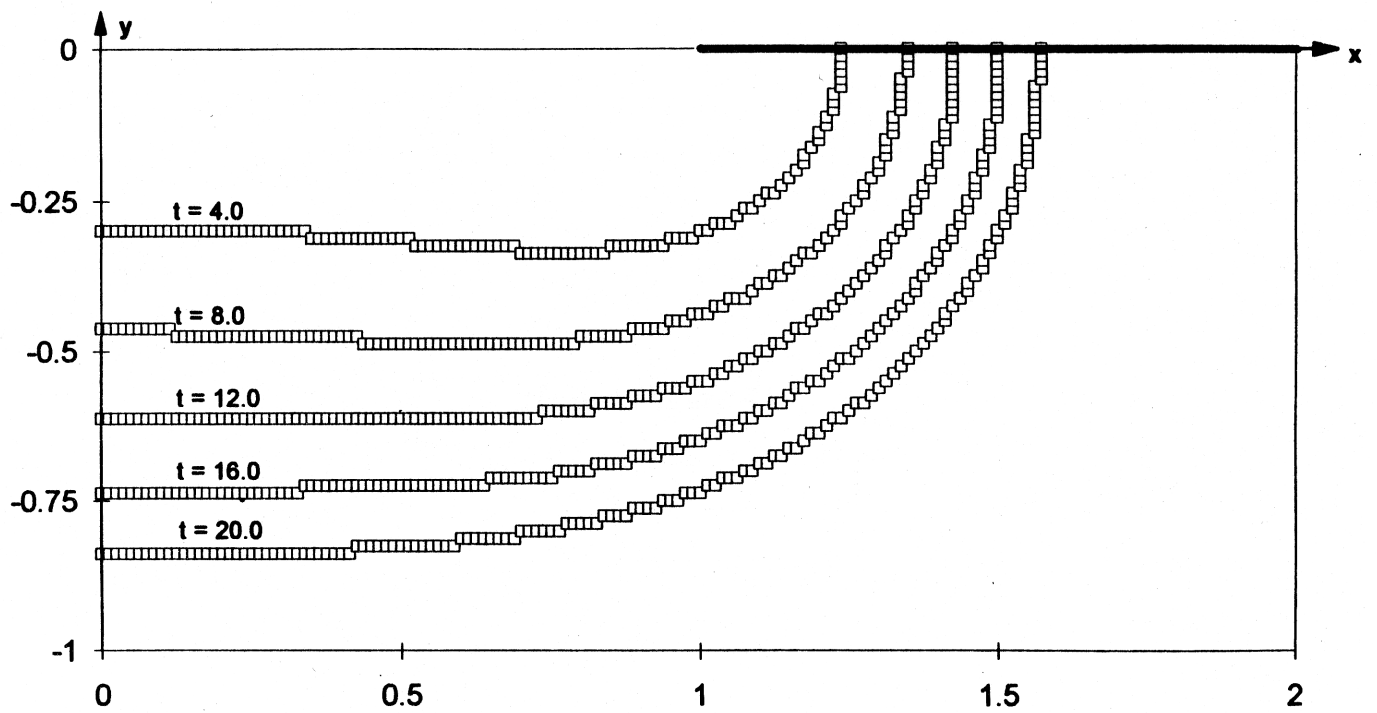


Figure 4: Moving boundary at various times.